

Mathematical Morphology and Applications - End-Sem Exam

B.Math.(Hons.) Third Year

May 5, 2017

Instructions: There are **7** questions altogether. Marks corresponding to each question is indicated in bold. Answer as many as you can. Maximum score : **60 marks**. Maximum time : **3 hrs**.

- (1) (a) Define Serra's method to calculate the morphological median.
(b) Calculate the dual median of the definition above.
(c) Characterize the difference between the median and dual median.

[3+3+3]

- (2) Suppose X is a 2-D $n \times n$ binary image on a square grid. One can implement a dilation of X w.r.t $kB = B \oplus B \oplus \dots \oplus B$ (k times) or a multi-scale dilation of X w.r.t. B where $B = \{(-1, 0), (0, -1), (0, 0), (0, 1), (1, 0)\}$ using either Algorithm 1 k successive times or Algorithm 2 (see below).
- (a) Complete Algorithm 3 so that it has $\mathcal{O}(n^2)$ time complexity and $\mathcal{O}(1)$ space complexity. Also, justify the correctness of your algorithm.
- (b) Which among the following is more efficient: implementing Algorithm 1 k times or Algorithm 2? You will have to justify your answer by comparing the asymptotic time and space complexities (note that these expressions are functions of both n and k).

Algorithm 1 A Dilation Algorithm

```
1: procedure DILATE( $X, B$ )                                ▷ The dilation of  $X$  by  $B$ 
2:    $l \leftarrow X.shape[0]$  and  $b \leftarrow X.shape[1]$       ▷ Identify the length and width of the image  $X$ 
3:   for  $i < l$  do                                          ▷ Iterating over all the rows
4:     for  $j < b$  do                                          ▷ Iterating over all the columns
5:       if  $X[i][j] == 1$  then                                ▷ Lines 5-13: Marking the non-foreground pixels adjacent to
         foreground pixels
6:         if  $i > 0$  and  $X[i-1][j] = 0$  then
7:            $X[i-1][j] = 2$ 
8:         if  $j > 0$  and  $X[i][j-1] = 0$  then
9:            $X[i][j-1] = 2$ 
10:        if  $i < l-1$  and  $X[i+1][j] = 0$  then
11:           $X[i+1][j] = 2$ 
12:        if  $j < b-1$  and  $X[i][j+1] = 0$  then
13:           $X[i][j+1] = 2$ 
14:   for  $i < l$  do                                          ▷ Iterating over all the rows
15:     for  $j < b$  do                                          ▷ Iterating over all the columns
16:       if  $X[i][j] == 2$  then
17:          $X[i][j] = 1$ 
18:   return  $X$                                               ▷ Outputs the dilated image  $X$ 
```

[8+4]

Algorithm 2 A Multiscale Dilation

```
1: procedure MULTISCALEDILATE( $X, B, k$ ) ▷ The dilation of  $X$  by  $kB$ 
2:    $l \leftarrow X.shape[0]$  and  $b \leftarrow X.shape[1]$  ▷ Identify the length and width of the image  $X$ 
3:    $X = \text{TAXICAB}(X)$  ▷ Compute the distances of every pixel from a nearest foreground pixel
   w.r.t taxicab metric
4:   for  $i < l$  do ▷ Iterating over all the rows
5:     for  $j < b$  do ▷ Iterating over all the columns
6:       if  $X[i][j] \leq k$  then
7:          $X[i][j] = 1$ 
8:       else
9:          $X[i][j] = 0$ 
10:  return  $X$  ▷ Outputs the multi-scale dilated image  $X$ 
```

Algorithm 3 Distances from Nearest Foreground Pixel

```
1: procedure TAXICAB( $X$ ) ▷ The distances of every pixel in  $X$  from a nearest foreground pixel w.r.t.
   taxicab metric
2:   ...
3:   ...
4:   ...
5:  return  $X$  ▷ Outputs a 2-D array  $X$  containing the nearest distances from foreground
```

- (3) Let B be a convex set in $\mathcal{P}(\mathbb{R}^2)$. Define

$$\mathcal{L} = \{X \oplus B, \text{ for all } X \subseteq \mathbb{R}^2\}$$

Define the order relation as - $X \leq Y$ if and only if $X \subseteq Y$. Is (\mathcal{L}, \leq) a lattice? If yes, characterize the sets $X \wedge Y$ and $X \vee Y$ for any two sets $X, Y \in \mathcal{L}$.

[4+(2+2)]

- (4) Recall that the definition of dilation $\delta(\cdot)$ on $\mathcal{P}(\mathbb{R}^2)$ is given by

$$\delta_B(X) = \bigcup_{b \in B} X_b$$

Dilation can also be defined on a lattice (\mathcal{L}, \leq) as the increasing operator which preserves the supremum. i.e.,

$$\delta \left(\bigvee_i x_i \right) = \bigvee_i \delta(x_i)$$

Stating the minimum required assumptions, show the equivalence between these two definitions. What can be said about the erosion operator?

[6+3]

- (5) Let X be a disconnected compact set in \mathbb{R}^2 w.r.t. euclidean metric. Define the $SKIZ(\cdot)$ operator on X and assume that for every point $x \in SKIZ(X)$, there is an associated distance $d(x) \in \mathbb{R}^+$ which denotes the distance from x to the nearest point in X . Prove or disprove:

$$X^c = \bigcup_{x \in SKIZ(X)} B(x, d(x))$$

where $B(x, d)$ is the closed ball with center x and radius d .

[5]

- (6) (a) Define a white *ASF* (alternating sequential filter) and describe some cases where *ASF* is preferable over multi scale openings.
- (b) Can a white/black *ASF* and *Id* (identity) operators be compared? If yes state the relation and prove it. If not explain with the help of counter examples.

[(1+3)+(1+3)]

- (7) Indicate *True* or *False* for each of the following statements and provide reasons supporting your answers. You will have to prove the statement if you indicate *True* and provide a counter example if you indicate *False*. Correct answer carries 1 mark each and a valid reasoning carries 4 marks each.

- (a) Let $\mathcal{G}_1 = (V, E, W_1)$ be a vertex-weighted graph. Suppose $\mathcal{G}_2 = (V, E, W_2)$ is another vertex-weighted graph such that $W_2(x) = W_1(x) - 1$ where x is a destructible point of \mathcal{G}_1 and $W_2(y) = W_1(y)$ for every $y \in V \setminus \{x\}$. Let $M(\mathcal{G})$ denote the minima of \mathcal{G} (which is a sub-graph of \mathcal{G}). Assume that $M(\mathcal{G}_1)$ has at least two connected components. The pass value between a given pair of connected components of $M(\mathcal{G}_1)$ in \mathcal{G}_1 is same as the pass value between the same pair in \mathcal{G}_2 .
- (b) Suppose $\mathcal{G} = (V, E, W)$ is a vertex-weighted graph then topological watershed of \mathcal{G} is unique.
- (c) Suppose $\mathcal{G} = (V, E, W)$ be an edge-weighted graph. Let M denote the minima of the graph \mathcal{G} . A watershed cut i.e. a cut induced by a minimum spanning forest relative to M is a max-cut relative to M .

[3 × (1+4) = 15]